

GAMES FOR LEARNING

A SABOTAGE APPROACH

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OUTLINE

- 1 INTRODUCTION
- 2 LEARNING AS A SABOTAGE GAME
 - Sabotage Modal Logic
 - Learning in Sabotage Modal Logic
 - Complexity of Sabotage-Type Learning
- 3 NON-STRICT ALTERNATION
- 4 MORE ON LEARNING AND GAMES
- 5 CONCLUSIONS



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AIM

To highlight the *interactive* nature of the learning process by showing its relation with games.



OUR PERSPECTIVE ON LEARNING

- High-level analysis of inductive inference.
- Singling out a correct hypothesis from a range of possibilities.
- Many steps of “update” before the conclusion is reached.



SUCCESSFUL LEARNING

DEFINITION

Learner identifies Teacher's choice *in the limit* if after some finite number of guesses his choices stabilize on a correct hypothesis.

DEFINITION

Learner *finitely* identifies Teacher's choice if after some finite number of guesses he makes the right choice.



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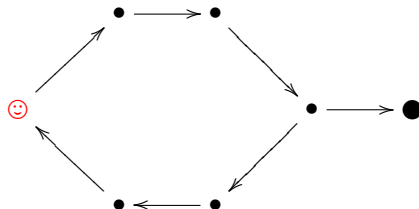
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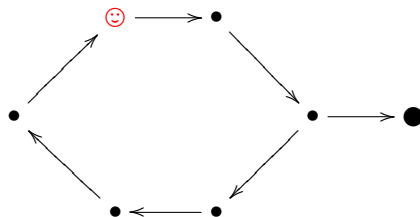
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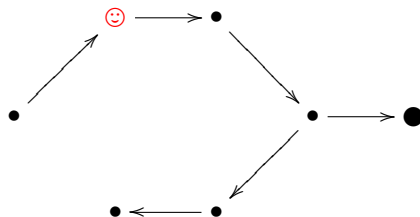
SABOTAGE GAME – EXAMPLE



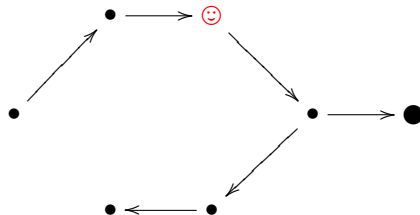
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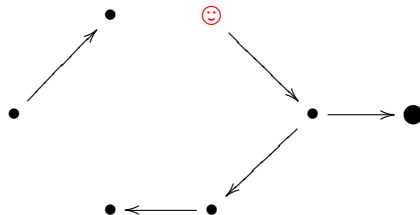
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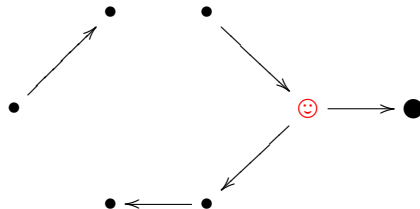
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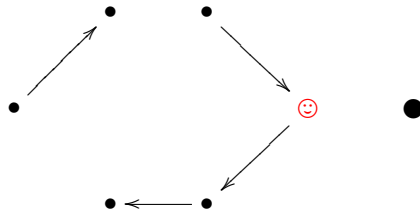
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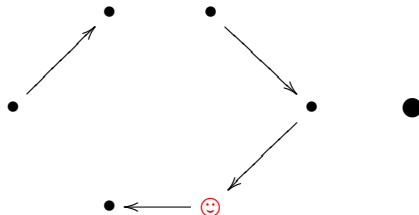
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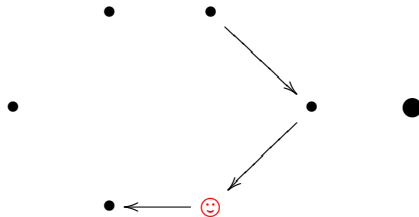
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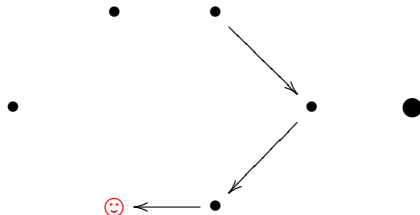
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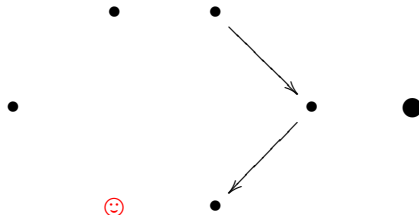
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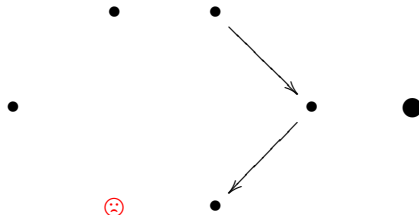
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SABOTAGE MODAL LOGIC

DEFINITION (SABOTAGE MODAL LANGUAGE)

$$\phi ::= p \mid \neg\phi \mid \phi \vee \phi \mid \Diamond_a\phi \mid \Diamond_a\phi$$

with $p \in \text{PROP}$ and $a \in \Sigma$ (finite).

$$\Diamond\phi := \bigvee_{a \in \Sigma} \Diamond_a\phi$$

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DEFINITION (SABOTAGE MODEL)

$M = \langle W, \{R_a \mid a \in \Sigma\}, Val \rangle$ where

$$W \neq \emptyset, \quad R_a \subseteq W \times W, \quad Val : \text{PROP} \rightarrow \mathcal{P}(W)$$



SABOTAGE MODAL LOGIC

DEFINITION (REMOVAL OPERATION)

Let $M = \langle W, \{R_a \mid a \in \Sigma\}, Val \rangle$ be a Sabotage Model.

$$M_{(u,v)}^a := \langle W, \{R_b \mid b \in \Sigma \setminus \{a\}\} \cup R_a \setminus \{(u, v)\}, Val \rangle$$

DEFINITION (SEMANTICS)

$$M, w \models \Diamond_a \phi \quad \text{iff} \quad \text{there is } (u, v) \in R_a \text{ s. t. } M_{(u,v)}^a, w \models \phi$$

THEOREM

Model checking of SML is PSPACE-complete.



SABOTAGE LEARNING GAME

DEFINITION

A *Sabotage Learning Game* is a Sabotage Game played between *Learner* and *Teacher* on a directed multi-graph with an initial vertex and a “goal” vertex.



VARIOUS SCENARIOS

Game	Winning Condition
SLG_{UE}	Learner wins iff he reaches the goal state, Teacher wins otherwise.
SLG_{HU}	Teacher wins iff Learner reaches the goal state, Learner wins otherwise.
SLG_{HE}	Both players win iff Learner reaches the goal state, Both lose otherwise.



CHARACTERIZATION RESULTS

THEOREM

Game	Existence of winning strategy	Winner
SLG_{UE}	$\gamma_0^{UE} := goal,$ $\gamma_{n+1}^{UE} := goal \vee \Diamond \Box \gamma_n^{UE}$	<i>Learner</i>
SLG_{HU}	$\gamma_0^{HU} := goal,$ $\gamma_{n+1}^{HU} := goal \vee (\Diamond \top \wedge (\Box \Diamond \gamma_n^{HU}))$	<i>Teacher</i>
SLG_{HE}	$\gamma_0^{HE} := goal,$ $\gamma_{n+1}^{HE} := goal \vee \Diamond \Diamond \gamma_n^{HE}$	<i>Both</i>



Complexity of Sabotage-Type Learning

COMPLEXITY OF SABOTAGE-TYPE LEARNING

Game	Winning Condition	Complexity
SLG_{UE}	Learner wins iff he reaches the goal state, Teacher wins otherwise	PSPACE-complete.
SLG_{HU}	Teacher wins iff Learner reaches the goal state, Learner wins otherwise.	PSPACE
SLG_{HE}	Both players win iff Learner reaches the goal state. Both lose otherwise.	NL-complete.



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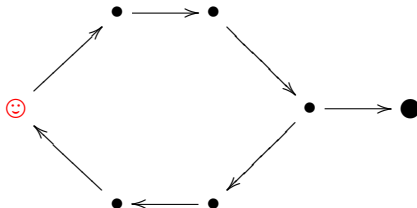


LOCAL VS GLOBAL MOVES

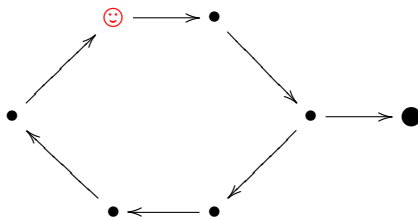
- Players moves are of a different nature:
 - Learner moves by *local* transitions.
 - Teacher moves by *globally* removing an edge.
- Teacher only needs to act when necessary.



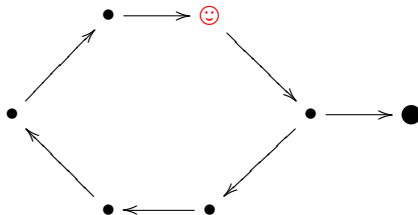
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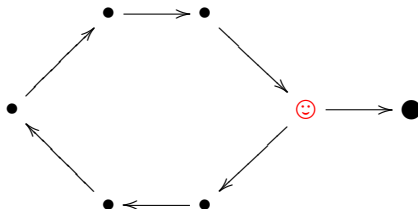
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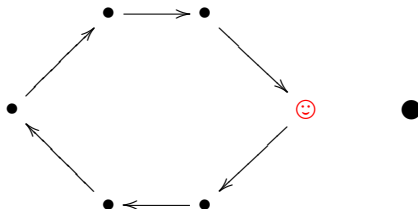
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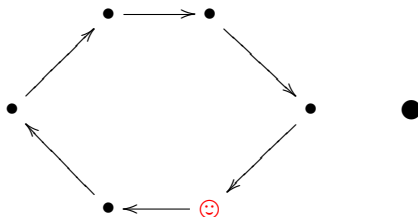
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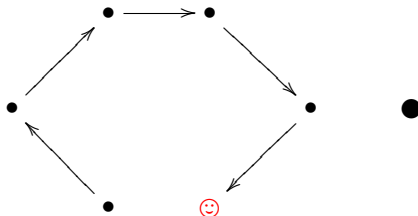
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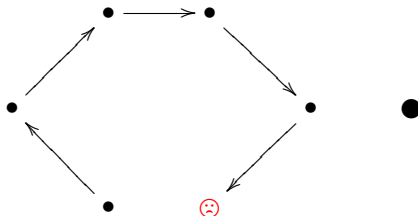
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NON-STRICT ALTERNATION – EXAMPLE



DIFFERENT GAME, SAME SCENARIOS

Game	Winning Condition
SLG_{UE}^*	Learner wins iff he reaches the goal state, Teacher wins otherwise.
SLG_{HU}^*	Teacher wins iff Learner reaches the goal state, Learner wins otherwise.
SLG_{HE}^*	Both players win iff Learner reaches the goal state, Both lose otherwise.



STRICT VS NON-STRICT ALTERNATION

THEOREM

- ① *Learner has a w.s. in SLG_{UE}^* iff he has a w.s. in SLG_{UE} .*
- ② *Teacher has a w.s. in SLG_{HU}^* iff she has a w.s. in SLG_{HU} .*
- ③ *Teacher and Learner have a joint w.s. in SLG_{HE}^* iff they have a joint w.s. in SLG_{HE} .*

COROLLARY

Formulas provided before characterize existence of a winning strategy in SLG_{UE}^* , SLG_{HU}^* and SLG_{HE}^* , resp.



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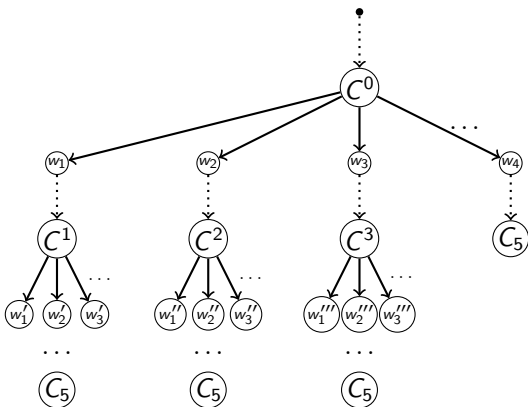


GAME THEORY AND LEARNING THEORY

- The use of GT in LT.
- The use of LT in Games.



THE GAME OF QUERIES AND COUNTEREXAMPLES



GAME THEORETIC APPROACH TO INDUCTIVE INFERENCE

- Epistemic status of the players, imperfect information, payoff characteristics.
- Choice for Learner
 - ① at each step the learner can choose from one or more procedures which are part of one algorithm;
 - ② in the beginning Learner can decide with which of the available algorithms he is going to proceed.



LEARNING THEORY IN GAMES

- Inductive inference games.
- Eleusis (identification in the limit), Zendo (queries and counterexamples).
- Complexity analysis of corresponding algorithms.
- Empirical work 1: difficulty vs. complexity of hidden rules.
- Empirical work 2: knowledge reports.



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SUMMARY AND CONCLUSION

AIM

Provide a high-level game-theoretical perspective on formal learning theory.

In particular, study strategic abilities, information flow and interaction.

SUMMARY

Game-theoretical approach to learning that accounts for different levels of cooperativeness between Learner and Teacher.



FURTHER WORK

- Identification in the limit (stable positions).
 - Epistemic and doxastic interpretation: operational, non-introspective knowledge.
 - Fixed-point logics for identification in the limit.
- GT approach to learning algorithms.
- Inductive inference games - theoretical and empirical account.



Dziękuję, Danke and Gracias

